

## Ford-Fulkerson's Maximum-Flow Algorithm

**Given:**

- (i) A connected directed graph  $G = (V, A)$ , two distinguished vertices called *source* ( $s$ ) and *sink* ( $t$ ),
- (ii) capacities  $u_{ij} \geq 0$  on each arc  $(i, j) \in A$ .

**Want:** A maximum value  $s \rightarrow t$  flow  $(\bar{x}_{ij})$  that observes capacities, and such that net flows into all vertices except  $s$  and  $t$  are zero.

**Initialize:** Pick any feasible flow  $\bar{x}$  (zero flow on each arc if no other is known).

Step 1 Define a residual graph,  $\hat{G} = (V, \hat{A})$  and capacities  $\hat{u}_{ij}$  with respect to the current flow,  $\bar{x}$  in the following way:

$$\hat{A} = \hat{A}_1 \cup \hat{A}_2$$

where

$$\hat{A}_1 = \{(i, j) : (i, j) \in A, \bar{x}_{ij} < u_{ij}\}$$

i.e. all the arcs in original graph, whose current flow values are not at their capacities; and

$$\hat{A}_2 = \{(j, i) : (i, j) \in A, \bar{x}_{ij} > 0\}$$

i.e. create a new arc  $(j, i)$  if arc  $(i, j)$  in original graph has positive current flow.

If  $(i, j) \in \hat{A}_1$  then assign **new capacity**  $\hat{u}_{ij} = u_{ij} - \bar{x}_{ij}$

If  $(j, i) \in \hat{A}_2$  then assign **new capacity**  $\hat{u}_{ji} = \bar{x}_{ij}$

Step 2 Find an augmenting chain in  $G$  i.e. find an  $s \rightarrow t$  directed path in  $\hat{G}$ .

If such directed path in  $\hat{G}$  exists then go to Step 3; otherwise go to Step 4.

Step 3 Augmentation or update:

Compute  $\delta \leftarrow \min\{\hat{u}_{ij} : (i, j) \text{ is in the directed } s \rightarrow t \text{ path in } \hat{G}\}$ .

Update

$$\bar{x}_{ij} \leftarrow \begin{cases} \bar{x}_{ij} + \delta & \text{if arc } (i, j) \text{ is forward on corresponding chain in } G \\ \bar{x}_{ij} - \delta & \text{if arc } (i, j) \text{ is reverse on corresponding chain in } G \\ \bar{x}_{ij} & \text{otherwise} \end{cases}$$

Repeat Step 1

Step 4 The current flow is optimal.

Define  $S \leftarrow \{v \in V : v \text{ is reachable from } s \text{ in } \hat{G}\}$ . STOP.

**Recall:** A vertex  $u$  is *reachable from* a vertex  $w$  in a directed graph if there is a directed path from  $w$  to  $u$  in the same directed graph.